

30 November 1983

STATEMENT OF SOME CURRENT RESULTS ABOUT THREE-WAY ARRAYS.

Joseph B Kruskal

max rank( $I \times J \times K$ ) means the maximum possible rank of an  $I \times J \times K$  array.  
If  $I \leq J \leq K$ , max rank satisfies

$$IJK / (I+J+K-2) \leq \text{max rank}(I \times J \times K) \leq IJ,$$

and some exact values are:

2 x 2 x K: for K=(1,2,3, >=4),	max rank=(2,3,3,4);
2 x 3 x K: for K=(1,2,3,4,5, >=6),	max rank=(2,3,4,5,5,6);
2 x 4 x K: for K=(1,2,3,4, 7, >=8),	max rank=(2,4,5,6, 7,8);
3 x 3 x K: for K=(1,2,3, >=9),	max rank=(3,4,5, 9).

The geometrical arrangement of arrays of different rank within the 8-dimensional space of  $2 \times 2 \times 2$  arrays has been worked out in considerable detail.

*Most difficult cases are circled.*

October 26, 1983  
Wednesday, 1:20pm

MH Room 6H-404

## RANK AND GEOMETRY OF THREE-DIMENSIONAL MATRICES

Joseph B Kruskal (Dept 11214)

The concept of rank extends naturally to 3-dimensional matrices, and the extension has applications to complexity theory and statistics. Calculating rank is difficult even for tiny matrices. We present new results up to  $3 \times 3 \times 3$ . Max  $2 \times 2 \times 2$  rank is 3; surprisingly, both rank 2 and rank 3 occur with positive measure. Their geometrical arrangement is described.

RANK & GEOMETRY

of THREE-DIMENSIONAL

MATRICES

Joseph B Kruskal

26 Oct 1983

Rank of a matrix is a major mathematical concept. There is a natural generalization to  $n$ -dim matrices or tensors having rank [in different sense]  $n$ .

It is so natural that it has arisen at least 4 times, apparently independently:

- ① computational complexity - Strassen (1968) and many subsequent papers
- ② psychometrics/statistics - Carroll & Chang ('70), Harshman (1970), etc.
- ③ abstract algebra - Watkins (1976)
- ④ differential geometry? - [reference mislaid]

I will present it as a concept of pure mathematics.

The most elegant approach defines rank for sets of  $n$ -dim. matrices.

All results for real numbers  $\mathbb{R}$  only.

First I will define dim and rank in a formal way. Then I will connect them to matrix rank, etc

"dim" has 2 meanings		"rank" has 2 meanings
- as above		- as above
- as in " <del>N-dim matrix</del> "		- as in " <del>rank tensor</del> "

Use N-way array or N-way tensor instead

N-way array has elements  $X_{i_1 \dots i_N}$

Vector is 1-way array. Matrix is 2-way array.

(  $\mathbb{R}^I$  = vectors (I-tuples) = I arrays

$\mathbb{R}^{I \times J}$  =  $I \times J$  matrices =  $I \times J$  arrays

(  $\mathbb{R}^{I \times J \times K}$  =  $I \times J \times K$  arrays

Size of array is  $I$  or  $(I, J)$

or  $(I, J, K)$  etc

Arrays (of fixed size) form a vector space.

$\mathcal{V}$  is a set of arrays (of same size)

Def Linear span of  $\mathcal{V}$  is set of all linear combinations of  $\mathcal{V}$

= linear space (i.e., vector space) generated by  $\mathcal{V}$

= smallest linear space containing  $\mathcal{V}$ .

If  $\mathcal{V}$  is a linear space, then

$\mathcal{V}$  = linear span of  $\mathcal{V}$

Def.  $\dim(\mathcal{V})$  = min # of arrays whose linear span contains  $\mathcal{V}$ .

If  $\mathcal{V}$  is a linear space, then  $\dim(\mathcal{V})$  agrees with usual meaning.

If  $\mathcal{V}$  consists of  $n$  indep elts,  $\dim(\mathcal{V}) = n$

If matrix  $m \neq 0$ , array  $X \neq 0$ ,

$\dim(m) = 1$ ,  $\dim(X) = 1$ .

$\dim(\mathcal{V}) = \dim(\text{linear span of } \mathcal{V})$

Def An  $N$ -way array is an outer product array if it is the outer product of  $N$  vectors.

Meaning

Meaning of -

Outer product matrix:  $m_{--} = a_{-} \otimes b_{-}$ ,

$$m = \begin{bmatrix} a_1 \\ \vdots \\ \vdots \\ \vdots \end{bmatrix} \begin{bmatrix} b_1 & \dots \end{bmatrix}, \quad m_{ij} = a_i b_j$$

Dyad, dyadic matrix, "multiplication table",  
rank-1 matrix, ...

Outer product 3-way array  $X_{---} = a_{-} \otimes b_{-} \otimes c_{-}$

$$X_{ijk} = a_i b_j c_k$$

Triad, triadic array, rank-1 array

Outer product vector: ANY vector!  
Vacuous restriction.

$\mathcal{V}$  set of arrays (of same size)

Def  $\text{rank}(\mathcal{V}) = \min \#$  of outer product arrays whose linear span contains  $\mathcal{V}$ .

Compare to def of "dim"

Fact:  $\text{rank}(\mathcal{V}) \geq \text{dim}(\mathcal{V})$

Fact:  $\text{rank}(\mathcal{V}) = \text{rank}(\text{Linear span of } \mathcal{V})$

If  $\mathcal{V}$  consists of single matrix, this def agrees with classical meaning of rank.

If  $\mathcal{V}$  consists of vectors (1-way arrays),

$$\text{rank}(\mathcal{V}) = \text{dim}(\mathcal{V}),$$

because all vectors are "outer product vectors".

If  $\mathcal{V}$  is empty or  $\mathcal{V}$  consists of 0 array,  $\text{rank}(\mathcal{V}) = 0$ , and conversely.

Suppose  $X$  is  $I \times J \times 1$  array, &  $X_{ij1} = M_{ij}$ .

Then  $\text{rank}(X) = \text{rank}(M)$ ; and similar result holds in general.



# SOME EXAMPLES

$\mathcal{V}$	$\dim \leq$	$\text{rank}$	generating set of outer product mats
$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$	1	1	$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$
$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	1	2	$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$
$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$	2	2	$\begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}, \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix}$
$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 \\ -1 & 0 \end{bmatrix}$	2	3	$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 1 \\ -1 & -1 \end{bmatrix}$
Any 2 2x2 matrices	$\leq 2$	$\leq 3$	
Any 3 ...	$\leq 3$	$\leq 3$ same!	
Any 4 or more ...	$\leq 4$	$\leq 4$	$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$

Classical: Matrix also has row-rank, col-rank

$$\text{row-rank}(m) \stackrel{\text{def}}{=} \dim(\text{rows of } m)$$

$$\text{col-rank}(m) \stackrel{\text{def}}{=} \dim(\text{cols of } m)$$

Lemma: row-rank = col-rank = rank  
(over any field; They may differ over a ring).  $\square$

We incorporate this as follows.

Def: If  $X$  is  $N$ -way array of size  $(I_1, \dots, I_N)$   
and elts  $X_{i_1, \dots, i_N}$ , then the  $I_n$

$n$ -subarrays of  $X$  (for  $n=1$  to  $N$ ) are

$$X_{\dots i_n \dots} \quad \text{with } i_n = 1 \text{ to } I_n.$$

1-subarrays of matrix  $m$  are rows  $m_{i-}$

2- " " " " " " cols  $m_{-j}$

1-subarrays of 3-way array  $X$  are 1-slabs  $X_{i--}$

2- " " " " " " 2-slabs  $X_{-j-}$

3- " " " " " " 3-slabs  $X_{--k}$

Def  $n$ -rank( $X$ ) = rank (the  $n$ -subarrays of  $X$ ).

Note: We have not defined for sets of arr.

For matrix  $m$ ,

$$\begin{aligned} 1\text{-rank}(m) &= \text{rank}(1\text{-subarrays}) = \text{rank}(\text{rows}) \\ &= \dim(\text{rows}) = \text{row-rank} \end{aligned}$$

$$2\text{-rank}(m) = \text{col-rank}(m).$$

Lemma: For any  $N$ -way  $X$ ,

$$1\text{-rank} = 2\text{-rank} = \dots = N\text{-rank} = \text{rank}.$$

Could define  $n$ -rank for sets of arrays,  
but this result would not hold.

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Def:  $n$ -dim( $X$ ) = dim (the  $n$ -subarrays of  $X$ )

$$\text{For matrix } m, \quad \text{rank} = 1\text{-rank} = 2\text{-rank}$$

$\quad \quad \quad \parallel \quad \quad \quad \parallel$   
 $\quad \quad \quad 1\text{-dim} \quad \quad \quad 2\text{-dim}$

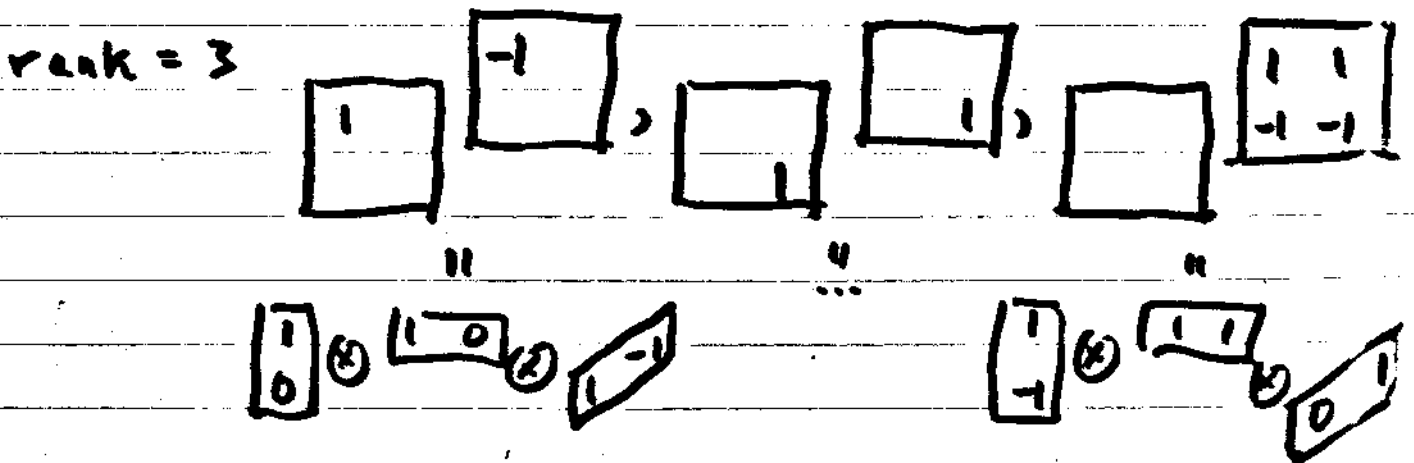
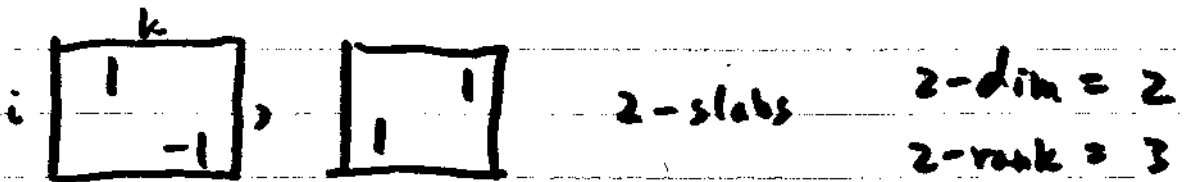
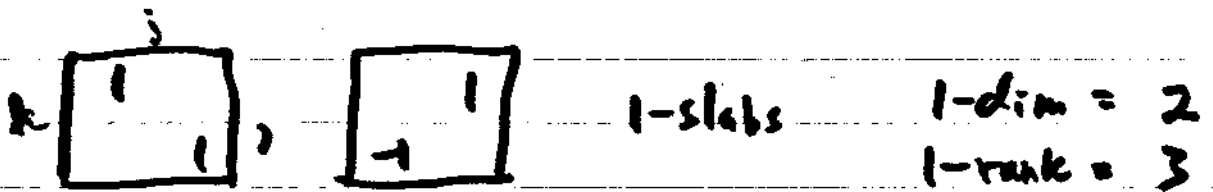
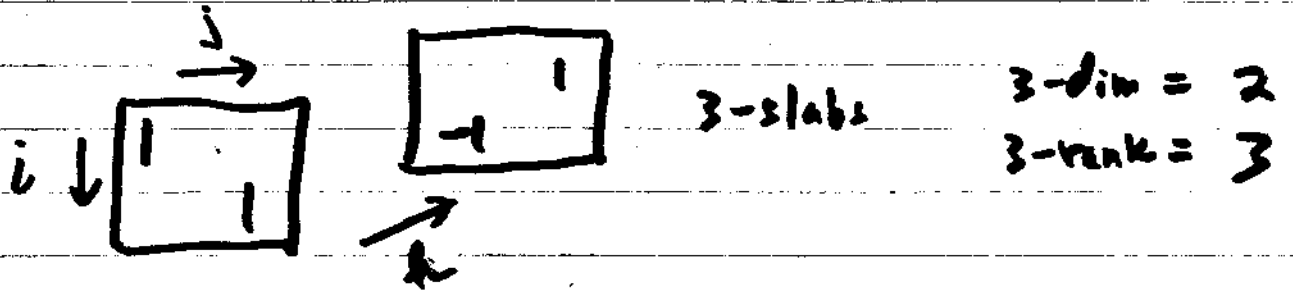
For 3-way array  $X$ ,

$$\text{rank} = 1\text{-rank} = 2\text{-rank} = 3\text{-rank}$$

$\quad \quad \quad \vee \parallel \quad \quad \quad \vee \parallel \quad \quad \quad \vee \parallel$   
 $\quad \quad \quad 1\text{-dim} \quad \quad \quad 2\text{-dim} \quad \quad \quad 3\text{-dim}.$

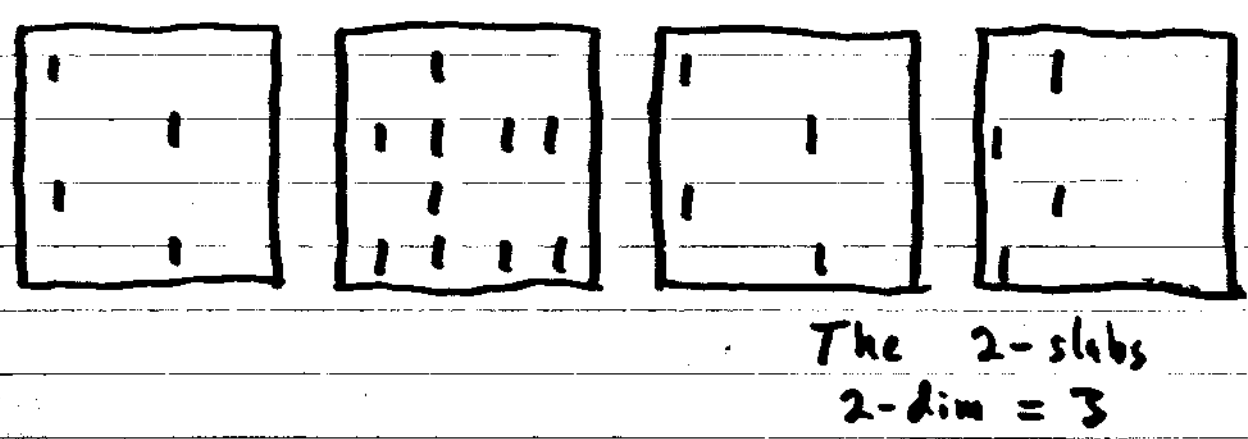
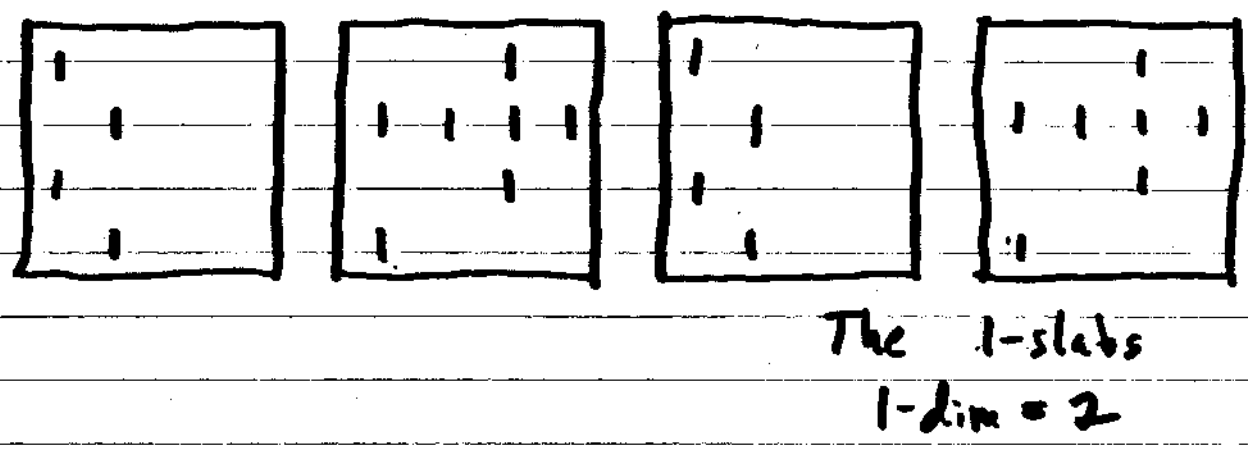
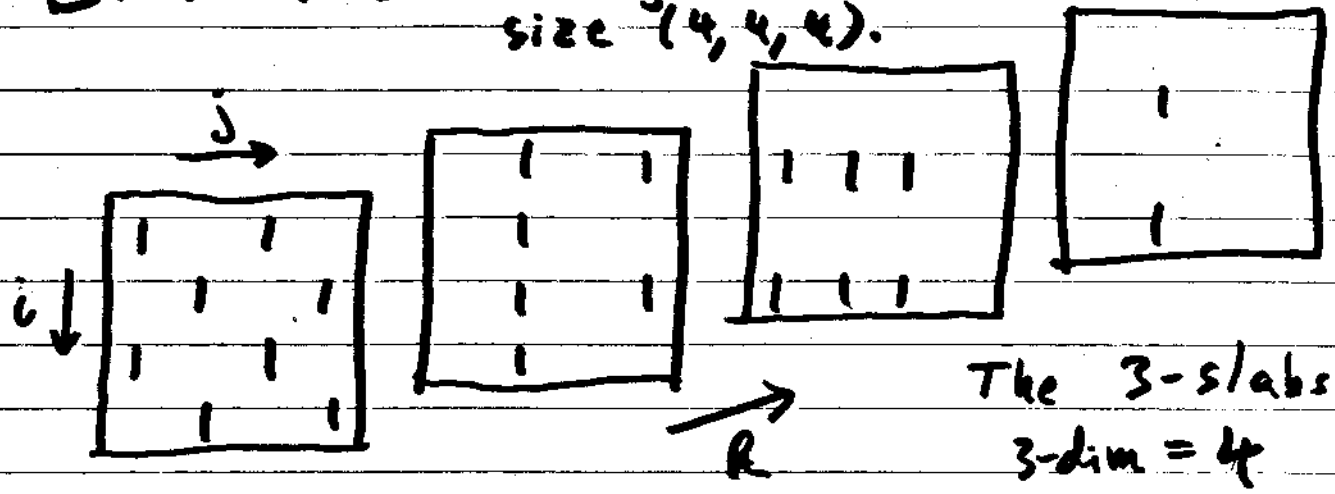
# EXAMPLE - Array of size (2, 2, 2)

Omitted elements = 0



rank	=	1-rank	=	2-rank	=	3-rank
3		3		3		3
		VII		VII		VII
		1-dim		2-dim		3-dim
		2		2		2

EXAMPLE - Array of size (4, 4, 4).



$$\text{Rank} = \begin{matrix} 1\text{-rank} & = & 2\text{-rank} & = & 3\text{-rank} \\ 4 & & 4 & & 4 \\ \text{VII} & & \text{VII} & & \text{VII} \\ 1\text{-dim} & & 2\text{-dim} & & 3\text{-dim} \\ 2 & & 3 & & 4 \end{matrix}$$

EXAMPLE - from Strassen algorithm

$$\begin{bmatrix} u_1 & u_2 \\ u_3 & u_4 \end{bmatrix} \begin{bmatrix} v_1 & v_2 \\ v_3 & v_4 \end{bmatrix} = \begin{bmatrix} w_1 & w_2 \\ w_3 & w_4 \end{bmatrix}$$

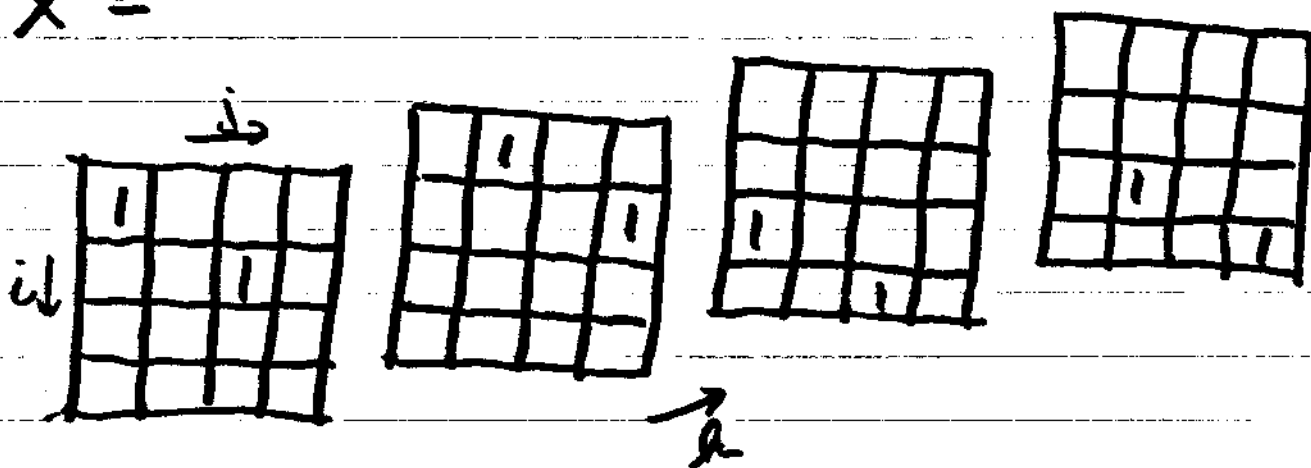
Then  $w_1 = u_1 v_1 + u_2 v_3$   
 $w_2 = \dots$ , etc

Define  $X_{ijk}$  to be the coeffs in

$$w_k = \sum_i \sum_j X_{ijk} u_i v_j$$

Then

$$X =$$



$$1\text{-dim} = 2\text{-dim} = 3\text{-dim} = 4$$

$$\text{rank} = 7$$

$\text{rank}(M) = \min \#$  of outer product matrices  
whose linear span contains  $M$

$= \min R$  for which it is possible to write

$$m_{ij} = \sum_{r=1}^R a_{ir} b_{jr} \quad \left\{ \begin{array}{l} \text{Each } r \text{ is one} \\ \text{outer product} \end{array} \right.$$

$$\text{or } M = \sum_I a_R b_J^T$$

$= \min R$  for which  $\sum_I M = \sum_I a_R b_J^T$  possible

$\text{rank}(X) = \min R$  for which it is possible to  
write

$$X_{ijk} = \sum_{r=1}^R a_{ir} b_{jr} c_{kr} \quad \text{"Triple product"}$$

$$\text{or } X_{(I,J,K)} = \left[ \sum_I a_R, \sum_J b_R, \sum_K c_R \right]$$

Decomposition means  $M = a b^T$  or

$$X = [a, b, c]$$

Rank & uniqueness - Kruskal (1977)

Comment - Two applications areas  
concerned with 3-way arrays.

Sometimes rank ( $X$ ) is central.

Sometimes decomposition is central.

Statistical PARAFAC model is based  
on least-squares-fitting by arrays  
of bounded rank.

Current work grew out of questions  
this stimulated. -- gave talk here in June.

Though this concept of rank goes back  
a few years, & can be found several  
places in the literature, this  
concise unified definition is new.



# SOME PREVIOUS RESULTS ON RANK

For  $I \times J \times K$  array:

$$\left. \begin{array}{l} 1\text{-dim} \\ 2\text{-dim} \\ 3\text{-dim} \end{array} \right\} \leq \text{rank} \leq \begin{cases} (1\text{-dim})(2\text{-dim}) \leq IJ \\ (1\text{-dim})(3\text{-dim}) \leq IK \\ (2\text{-dim})(3\text{-dim}) \leq JK \end{cases}$$

$$\text{rank} \leq \begin{cases} \# \text{ of non-zero 1-lines } X_{-jk} \\ \# \text{ " " " 2-lines } X_{i-k} \\ \# \text{ " " " 3-lines } X_{ij-} \end{cases}$$

Will need in proof later

[Frobenius (1911):  $\text{rank}(u^T v) + \text{rank}(v^T w) - \text{rank}(u^T w) \leq \text{rank}(u)$ ]

Defines: matrix  $\oplus$ , ARRAY:  $(u \oplus, X)_{ijk} = \sum_{\ell} u_{i\ell} X_{\ell jk}$

$$(u, v, w) \oplus X = u \oplus (v \oplus (w \oplus, X))$$

$$1\text{-dim}(u \oplus, X) + \text{rank}(v \oplus, w \oplus, X)$$

$$- 1\text{-dim}((u, v, w) \oplus X) \leq \text{rank}(X)$$

If 1-slabs of  $X$  are indep., then

$$\min_{\substack{\text{all } u \text{ with} \\ \text{lin. indep. rows}}} \text{rank}(u \oplus, X) + 1\text{-dim}(X) - 1\text{-dim}(u) \leq \text{rank}(X)$$

etc., etc., etc.

Lots of results - but lots of holes

Strassen array: rank 7; general results 6-8

HOLES ARE LARGE

See soon we can prove existence  
of high rank arrays but have very  
little ability to prove that individual  
arrays have that rank.

## SOME MOTIVATING QUESTIONS

What is max possible rank for  $I \times J \times K$  array?

How are arrays of each rank arranged geometrically in  $IJK$ -space?

Methods for calculating rank, or at least bounds on rank (for given array)?

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## Counting Parameters (rigorous argument)

$I \times J \times K$  array has  $IJK$  d.f.

Triple-product provides mapping

$$(A, B, C) \rightarrow [A, B, C]$$

$$\text{from } \mathbb{R}^{I \times R} \oplus \mathbb{R}^{J \times R} \oplus \mathbb{R}^{K \times R} \rightarrow \mathbb{R}^{I \times J \times K}$$

Since this mapping is smooth, it cannot be space-filling unless

$$(IR + JR + KR) \equiv (I + J + K)R \geq IJK.$$

Thus

$$\text{max rank} \geq \text{min space-filling rank} \geq \frac{IJK}{I + J + K}.$$

Improvement: If  $\Lambda, M, N$  are diagonal, and  $\Lambda M N = I$ , then

$$[A\Lambda, BM, CN] = [A, B, C].$$

$(\Lambda, M, N)$  has  $3R$  parameters; but only  $2R$  d.f. Then the mapping cannot be space-filling unless

$$(I + J + K)R - 2R \equiv (I + J + K - 2)R \geq IJK.$$

Thus

$$\text{min space-filling rank} \geq \frac{IJK}{I + J + K - 2}$$

For  $I \times I \times I$  arrays

$$\text{max rank} \geq \text{min space-filling rank} \geq \frac{I^3}{3I-2}$$

$I$	$I^3 / (3I-2)$
1	1
2	2
3	3.9 $\rightarrow$ 4
4	6.4 $\rightarrow$ 7 (compare Strassen array)
5	9.6 $\rightarrow$ 10

$I \approx \frac{1}{3} I^2$

For  $I \times I \times K$  max rank  $\geq \frac{I^2 K}{2I+K-2}$

$I=4$

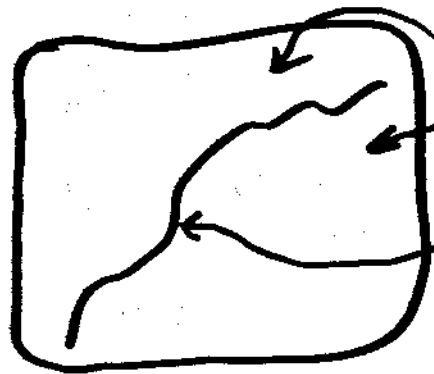
$K$	$I^2 K / (2I+K-2)$
1	2.3 $\rightarrow$ 3 [4]
2	4.0 $\rightarrow$ 4 ?
3	5.3 $\rightarrow$ 6 ?
4	6.4 $\rightarrow$ 7
5	7.3 $\rightarrow$ 8
6	8.0 $\rightarrow$ 8
7	8.6 $\rightarrow$ 9

$I=5$

$K$	$I^2 K / (2I+K-2)$
1	5.0 $\rightarrow$ 5 [5]
2	
3	
4	
5	9.6 $\rightarrow$ 10
6	
7	

# RANK OF $N \times N$ MATRICES (ILLUSTRATION)

$M_N$   
 $N^2$ -dim



Rank  $N$

Rank  $\leq N-1$

determinant = 0

Rank  $N$  "space-filling"  $N^2$ -dim

Lebesgue measure (= generalized volume)  
positive

Rank  $\leq N-1$  NOT space-filling,  $(N^2-1)$ -dim  
curved manifold  
measure 0

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Rank is  $N$  except for a set of  
measure 0

Rank is  $N$  almost everywhere

# ILLUSTRATION RANK OF $N \times N$ MATRICES

$\mathcal{M} =$  set of all  $N \times N$  matrices

	<u>Rank</u>	<u>Dimension</u>	<u>Lebesgue measure</u>
$\mathcal{M}_N = \mathcal{M}$	$\leq N$	$N^2$	<u>positive</u>

←  
[area,  
volume,  
⋮]

∪			
$\mathcal{M}_{N-1}$	<u>det=0</u> $\leq N-1$	$N^2-1$	0

∪			
$\mathcal{M}_{N-2}$	$\leq N-2$	$N^2-4$	0

∪			
$\mathcal{M}_{N-3}$	$\leq N-3$	$N^2-9$	0

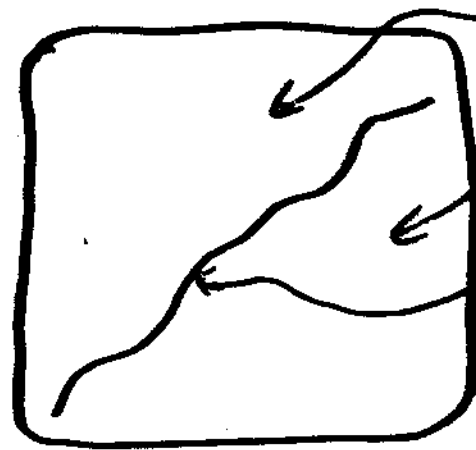
∪			
$\mathcal{M}_1$	<u>dyads</u> $\leq 1$	$2N-1$	0

∪			
$\mathcal{M}_0$	<u>zero matrix</u> 0	0	0

Each  $\mathcal{M}_n$  is curved manifold

# RANK OF $2 \times 2 \times 2$ ARRAYS

a  
8-dim



Rank 3 ↗ Different  
Rank 2 ↖  
Boundary: Rank = ?  
Can you guess?

- Rank 3      space-filling      8-dim
- Rank 2      space-filling      8-dim
- Boundary    NOT space-filling,    7-dim

## Fourth-degree homogeneous polynomial (12 terms)

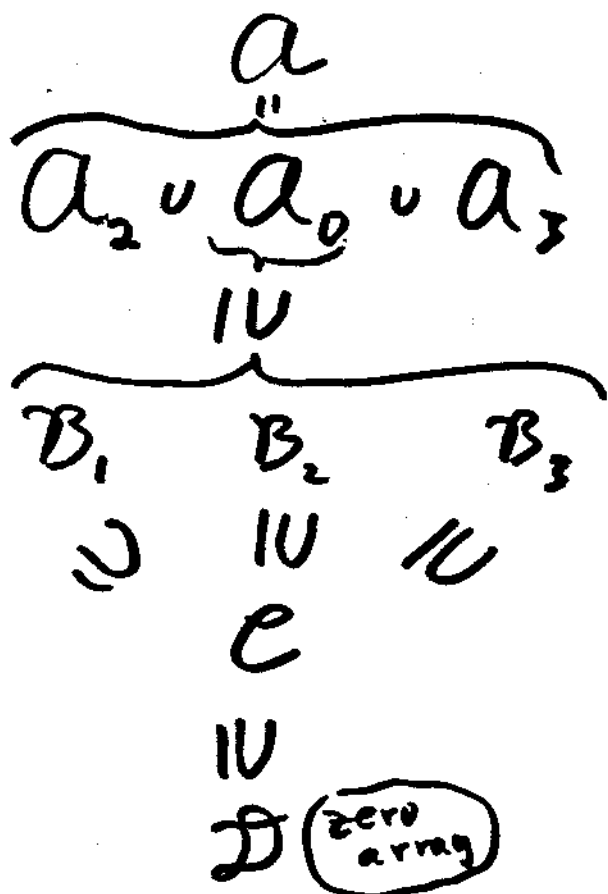
- positive      rank 3
- 0              boundary
- negative      rank 2

(akin to  
determinant  
of matrix)



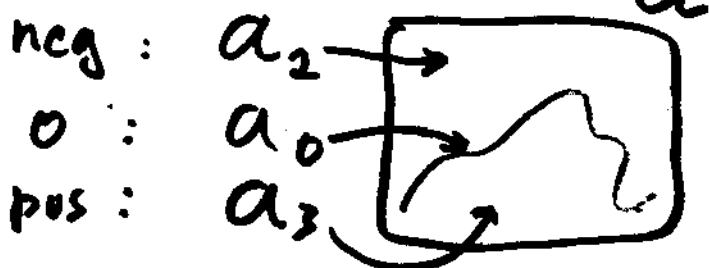
# RANK OF $2 \times 2 \times 2$ ARRAYS

$\mathcal{A}$  = set of all arrays



<u>Rank</u>	<u>Dim</u>	<u>Leb. meas</u>
	8	pos
$2 \leq 3 \leq 3$	8 7 8	pos 0 10s
$\leq 2 \leq 2 \leq 2$	5 5 5	0 0 0
$\leq 1$	4	0
0	0	0

gener. "determinant"  
(4th deg polynomial):



$$B_1: A_{1\dots} \sim A_{2\dots}$$

$$C = B_1 \cap B_2 = B_1 \cap B_3 = B_2 \cap B_3$$

= outer product arrays

ARRAY SIZE

RANK DISTRIBUTION

$1 \times J \times (K \geq J)$

0 1 2 ... J-1  $\boxed{J}$

$2 \times 2 \times 2$

0 1  $\boxed{2}$   $\boxed{3}$

$2 \times 2 \times 3$

0 1 2  $\boxed{3}$

$2 \times 2 \times (K \geq 4)$

0 1 2 3  $\boxed{4}$

$2 \times 3 \times 3$

0 1 2  $\boxed{3}$   $\boxed{4}$  ?

$2 \times 3 \times 4$

0 1 2 3  $\boxed{4}$  .

$2 \times 3 \times (K \geq 6)$

0 1 2 3 4 5  $\boxed{6}$

$3 \times 3 \times 3$

0 1 2 3  $\boxed{4}$   $\boxed{5}$   ~~$\boxed{6}$~~

Probably  
not  
possible

# A FEW PROOFS

Thm:  $\text{Rank}(2 \times 2 \times 2) \leq 3$ .

Proof: Since  $\text{rank} = 1 - \text{rank}$ , prove that any set of 2  $2 \times 2$  matrices has  $\text{rank} \leq 3$ .

Let  $\mathcal{V}$  be linear span of the set.

Know  $\dim(\mathcal{V}) \leq 2$ . Prove:  $\text{rank}(\mathcal{V}) \leq 3$

Case I:  $\dim(\mathcal{V}) \leq 1$ : Then  $\text{rank}(\mathcal{V}) =$   
rank of single  $2 \times 2$  matrix  $\leq 2$ .

Case II:  $\dim(\mathcal{V}) = 2$ , every matrix in  $\mathcal{V}$   
has  $\text{rank} \leq 1$ : Then any generating set  
consists of 2 outer product matrices, so  
 $\text{rank}(\mathcal{V}) \leq 2$ .

Case III:  $\dim(\mathcal{V}) = 3$ ,  $\mathcal{V} \ni$  matrix of rank 2:

Let  $\{a, b\}$  be basis with  $\text{rank}(a) = 2$ .

Then  $a^{-1}\{a, b\} = \{a^0, a^{-1}b\} = \{a^0, c\}$   
 $= \left\{ \begin{bmatrix} 1 & \\ & 1 \end{bmatrix}, \begin{bmatrix} c_{11} & c_{12} \\ c_{21} & c_{22} \end{bmatrix} \right\}$  is a basis.

Let  $d$  be a dyad (outer product matrix) which agrees with  $c$  in positions

$(1,2)$  &  $(2,1)$ :  $d^1 = d = \begin{bmatrix} ? & c_{12} \\ c_{21} & ? \end{bmatrix}$

E.g.,  $d_{11} \neq 0$ ,  $d_{22} = \frac{c_{12} c_{21}}{d_{11}}$

Then  $\{d^1, d^2 = \begin{bmatrix} 1 & \\ & \end{bmatrix}, d^3 = \begin{bmatrix} & \\ & 1 \end{bmatrix}\}$

is a dyadic (outer product) generating set for  $V$ , because

$$d \equiv \begin{bmatrix} 1 & \\ & 1 \end{bmatrix} = \begin{bmatrix} 1 & \\ & \end{bmatrix} + \begin{bmatrix} & \\ & 1 \end{bmatrix}$$

$$c = \begin{bmatrix} d_{11} & d_{12} \\ d_{21} & d_{22} \end{bmatrix} + (d_{11} - c_{11}) \begin{bmatrix} 1 & \\ & \end{bmatrix} + (d_{22} - c_{22}) \begin{bmatrix} & \\ & 1 \end{bmatrix}$$

Thm:  $\text{rank}(2 \times 2 \times 3) = 3$  almost everywhere.

rank  $\geq 3$  a.e.: The 3  $2 \times 2$  slices are indep a.e., so rank  $\geq 3$  a.e..

rank  $(2 \times 2 \times 3) \leq 3$ :

Def: 2 dyads  $u_1 \otimes v_1$  and  $u_2 \otimes v_2$  are compatible if either  $u_1 \sim u_2$  or  $u_2 \sim v_2$ .

Lemma: 3 pairwise incompatible  $2 \times 2$  dyads are indep.

Proof Let the dyads be  $u_i \otimes v_i = \begin{bmatrix} u_{i1} \\ u_{i2} \end{bmatrix} \begin{bmatrix} v_{i1} & v_{i2} \end{bmatrix}$

Consider a possible dependence,

$$0 = \sum_i \alpha_i \begin{bmatrix} u_{i1} \\ u_{i2} \end{bmatrix} \begin{bmatrix} v_{i1} & v_{i2} \end{bmatrix} = \sum \begin{bmatrix} u_{i1} \\ u_{i2} \end{bmatrix} \alpha_i \begin{bmatrix} v_{i1} & v_{i2} \end{bmatrix}$$

$$= \begin{bmatrix} u_{11} & u_{21} & u_{31} \\ u_{12} & u_{22} & u_{32} \end{bmatrix} \begin{bmatrix} \alpha_1 & 0 & 0 \\ 0 & \alpha_2 & 0 \\ 0 & 0 & \alpha_3 \end{bmatrix} \begin{bmatrix} v_{11} & v_{12} \\ v_{21} & v_{22} \\ v_{31} & v_{32} \end{bmatrix}$$

$$\equiv \bar{U} \bar{\alpha} \bar{V}$$

Case 0: No  $\alpha_i \neq 0$ . The eqn. is not a linear dependence.

Case 1: Just one  $\alpha_i \neq 0$ , say  $\alpha_1$ . Then  $u_1 \otimes v_1 = 0$ , so either  $u_1 = 0$  or  $v_1 = 0$ , say  $u_1$ . Then  $u_1 \sim u_2$ , contrary to hypothesis.

Case 2: Just two  $\alpha_i \neq 0$ , say  $\alpha_1 \neq \alpha_2$ . Then eqn reduces to

$$0 = \begin{bmatrix} u_{11} & u_{21} \\ u_{12} & u_{22} \end{bmatrix} \begin{bmatrix} \alpha_1 & 0 \\ 0 & \alpha_2 \end{bmatrix} \begin{bmatrix} v_{11} & v_{12} \\ v_{21} & v_{22} \end{bmatrix} = \bar{u} \bar{\alpha} \bar{v}$$

Now  $\text{rank}(\bar{\alpha}) = 2$ , and also

$$\begin{cases} \text{rank}(\bar{u}) = 2 \\ \text{rank}(\bar{v}) = 2 \end{cases} \begin{array}{l} \text{by incompatibility} \\ \text{hypothesis.} \end{array}$$

Thus  $\text{rank}(\bar{u} \bar{\alpha} \bar{v}) = 2$ , so  $\bar{u} \bar{\alpha} \bar{v} \neq 0$

Case 3: All  $\alpha_i \neq 0$ . Then  $\text{rank}(\bar{\alpha}) = 3$ ,  $\text{rank}(\bar{u}) = \text{rank}(\bar{v}) = 2$  by incompatibility.

Frobenius Theorem yields

$$\begin{aligned} \text{rank}(\bar{u} \bar{\alpha} \bar{v}) &\geq \text{rank}(\bar{u} \bar{\alpha}) + \text{rank}(\bar{\alpha} \bar{v}) - \text{rank}(\bar{\alpha}) \\ &= 2 + 2 - 3 = 1, \end{aligned}$$

so  $\bar{u} \bar{\alpha} \bar{v} \neq 0$ . This proves lemma.

Proof that  $\text{rank}(2 \times 2 \times 3) \leq 3$ .

$\text{rank} \geq 3 - \text{rank}$ , so it is enough to show that any  $\mathcal{V}$  generated by 3  $2 \times 2$  matrices has  $\text{rank} \leq 3$ . Since  $\mathbb{R}^{2 \times 2}$  is 4-dim,

$\exists m \in \mathbb{R}^{2 \times 2}$  orthog to  $\mathcal{V}$ .

If  $a, b \in \mathbb{R}^{2 \times 2}$ , their dot product  
 $= \sum_{i=1}^2 \sum_{j=1}^2 a_{ij} b_{ij} = \text{tr}(a b')$ ,

where  $\text{tr}(c) = \text{trace}(c) = \sum c_{ii}$ .

Thus  $\mathcal{V} = \{a \mid \text{tr}(a m') = 0\}$ .

It will suffice to find 3 indep dyads in  $\mathcal{V}$ .

Case 1:  $m$  is not a dyad.

For any non-zero <sup>column</sup> vector  $u$ ,  $u' m' \neq 0$ ,

so  $\exists$  <sup>col vector</sup>  $v = v(u)$  (unique up to scale)  $\Rightarrow u' m' v = 0$ ;

and  $v a' \in \mathcal{V}$  because

$$\text{tr}(v a' m') = \text{tr}(u' m' v) = 0.$$

Pick any 3  $u_i$  which are pairwise indep.

Then  $u_i^t m$  are pairwise indep, so  $v_i$  are

also. By the lemma,  $v_i v_i^t$  are indep.

Case 2:  $m$  is a dyad,  $m = x \otimes y = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \begin{bmatrix} y_1 & y_2 \end{bmatrix}$

$$m^t = (xy^t)^t = yx^t.$$

$$\mathcal{V} = \{a \mid \text{tr}(am^t) = 0\} = \{a \mid \text{tr}(ayx^t) = 0\}.$$

Set of dyads in  $\mathcal{V}$

$$= \{uv^t \mid 0 = \text{tr}(uv^t yx^t) = \text{tr}(x^t u v^t y) \\ = x^t u v^t y = (x \cdot u)(v \cdot y)\}$$

Thus  $uv^t \in \mathcal{V}$  iff  $u$  orthog to  $x$  or  
 $v$  orthog to  $y$

Let  $x^\perp$  &  $y^\perp$  be orthog to  $x$  &  $y$ , respectively.

Then  $x^\perp \otimes y^\perp$ ,  $x \otimes y^\perp$ ,  $x^\perp \otimes y$   
are dyads in  $\mathcal{V}$ . Indep. follows by  
simple direct argument.



# TOOLS I NEED THAT MAY BE KNOWN

Geometry of matrices

e.g., in  $\mathbb{R}^{I \times J}$ , what can we say

about intersection of  $p$ -dim

subspaces &  $\mathcal{M}_g$  (matrices of

rank  $\leq g$ )

Algebraic geometry